

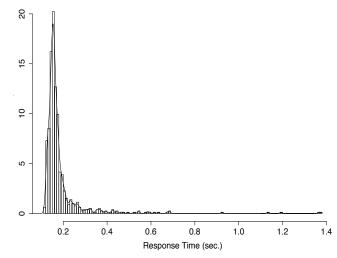




Heavy-tailed Distribution of Parallel I/O System Response Time

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Outline

Motivation

Response time sampling method

Analysis results of response time

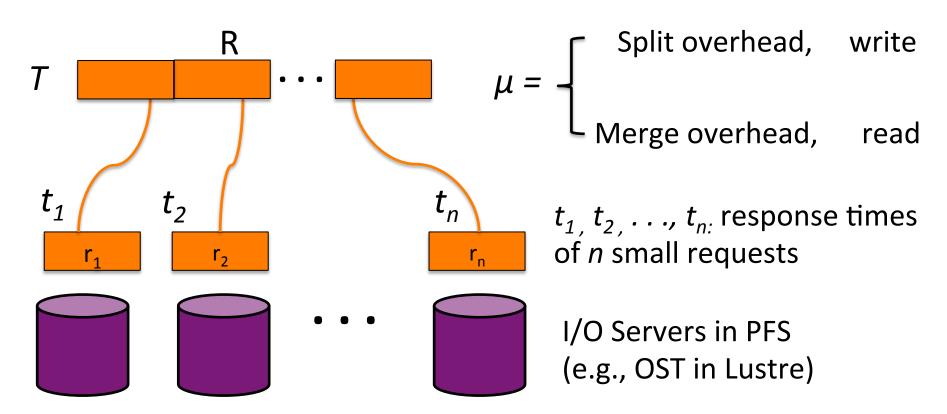
Estimating Response Time of I/O is Essential Element

- Data analysis query plan optimizing
 - Choose index or data organization with minimum read time
 - Scientific Data Services (SDS) framework,
 PostgresSQL, SciDB
- Data writing performance tuning
 - Select striping size, striping account, and other parameters to reduce write time
 - ExaHDF5, I/O Scheduler
- Simulator, Job Scheduler, Quality of service (QoS), etc.

Modeling Response Time for Parallel I/O

Response time of a single big file request R:

$$T = max(t_1, t_2, ..., t_n) + \mu$$



Simplifying Response Time Model

$$T = max(t_1, t_2, \ldots, t_n) + \mu$$

$$t_1$$

$$t_n$$

- Split/merge overhead μ is constant
- n small requests $\approx n$ sampling (i.i.d.) of n IO Servers
- $t_1, ..., t_n \approx n$ i.i.d. statistical variables
- Focus study on one (denoted by t) among t_1 , ..., t_n
 - -t: continuously distributed variable on $(0, +\infty)$

Applying Order Statistics to Estimate T

$$T = maximum(t_1, ..., t_n) + \mu$$
 $t : continuously distributed variable on $(0, +\infty)$
 $F_t(x) : distribution function of t$
 $f_t(x) = F_t'(x) : density function of t$$

- Step 1: Compute density function $f_{Yi}(y)$ with $F_t(x)$ and $f_t(x)$
 - Y_i : the **i-th** largest value $(t_1, t_2, ..., t_n)$ - $f_{Y_i}(y) = F(y)^{n-i} (1-F(y))^{n-i} f_t(y) n! / [(i-1)!(n-i)!]$ Order Statistics
- Step 2: Compute response time $T = Y_n$

Problem Statement

- What is the distribution function F(t) for the response time of each small file request?
 - Existing researches assume
 - Uniform Distribution
 - Normal Distribution
 - Are these assumptions true ?
 - If not, are there other distributions fitting better?

Our Method

 Sample the response time of two production storage systems

Analyze statistical properties of response time

Response Time Sampling Environments

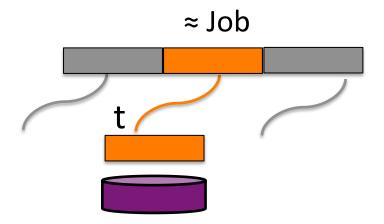
- Hopper and Edison at NERSC¹
 - 153K and 130K CPU cores, 1.28 PF and 2.39PF
 - 5000 registered users
 - 300 online active users on Edison
 - I/O Intensive jobs use Lustre
- Lustre file system
 - Cache on client and I/O server
 - Network latency
 - 1 ~ 143 OSTes

Computing Node /w Lustre Client Cache Network Router Cache Lustre OST

¹National Energy Research Scientific Computing Center https://www.nersc.gov/

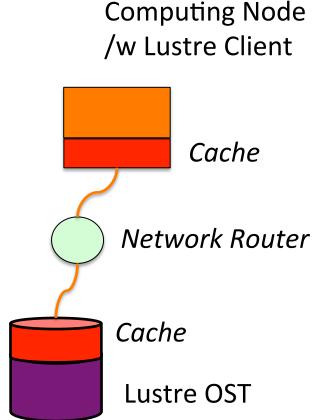
Sampling Method

- One job sampling one OST
 - A job ≈ A small file request
 - Measure time of reading and writing separately
 - Test different reading/writing sizes
 - 12 different sizes: 512KB, 1MB, 2MB, ..., 1024MB
 - Match request size and striping size



Sampling Method

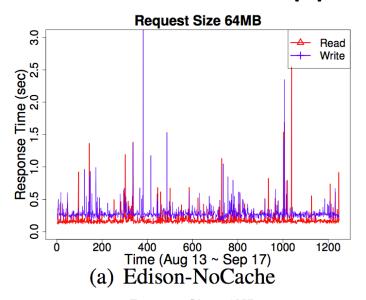
- Measure response time on computing node
 - network, disk, cache
- Cache Consideration
 - No Cache
 - clear cache by accessing memory sized data before sampling
 - call fsync() after write
 - Cache
 - High frequently sampling

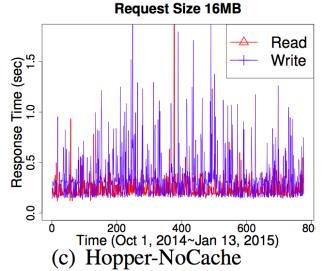


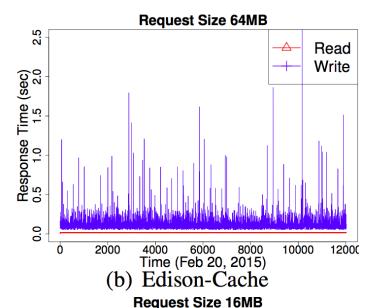
Sampling Results Statistics Overview

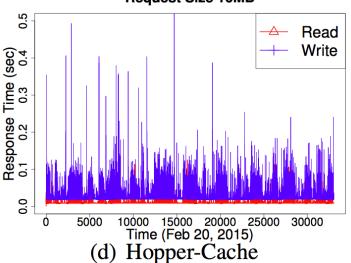
	Start Time	End Time	Days	# of Sampling	# of OSTs
Edison- No-Cache	08/13/2014	09/17/2014	35	14,977	12
Edison- Cache	02/20/2015	02/20/2015	1	927,691	12
Hopper- No-Cache	10/01/2014	01/13/2015	104	13,868	12
Hopper- Cache	02/20/2015	02/20/2015	1	1,581,364	12
		Summary	141	2,537,900	48

Variability of Raw Response Time for Edison and Hopper, Cache and No-Cache

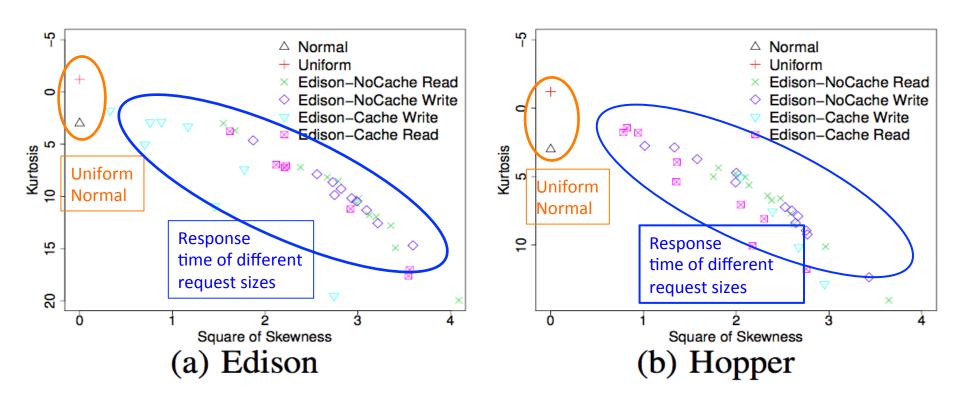








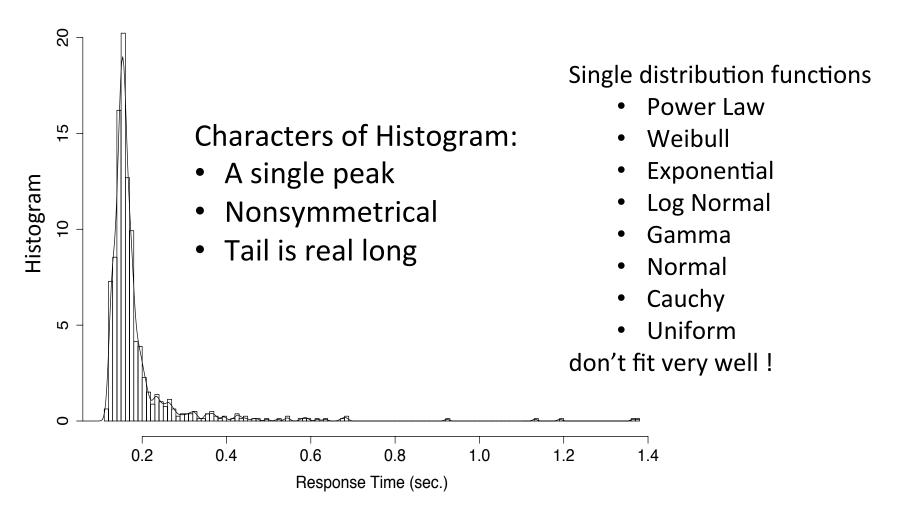
Ill-fit of Uniform or Normal Distribution



Metrics	Uniform	Normal
Kurtosis	- 1.2	3
Skewness	0	0

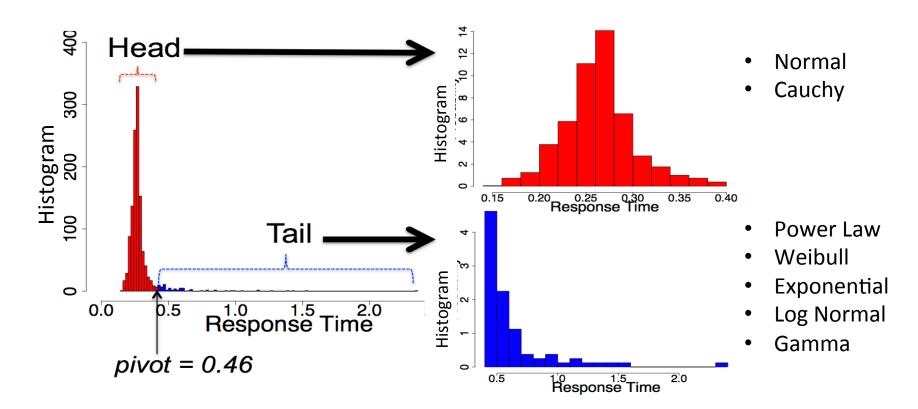
Ill-fit of Uniform, Normal, and Other Single Distribution Function

Read (Stripe Size: 64MB)



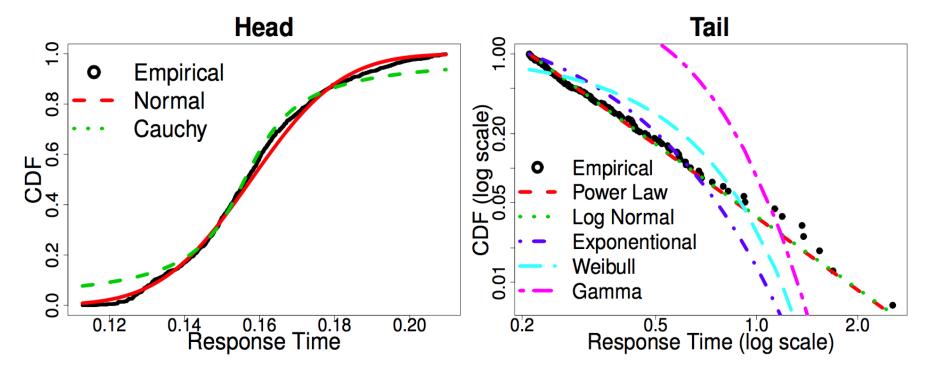
Exploring New Distributions

- Partition response time into Head and Tail
- Find the pivot
 - minimizing KS (Kolmogorov-Smirnov) distances



Fitting Results

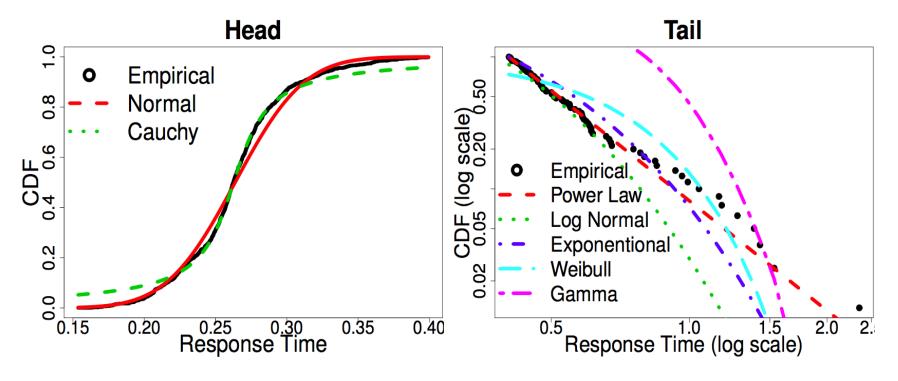
• Edison-NoCache, Read Response Time, 64MB



	Accuracy
Head Group	Normal > Cauchy
Tail Group	Power Law > Log Normal > Exponential > Weibull > Gamma

Fitting Results

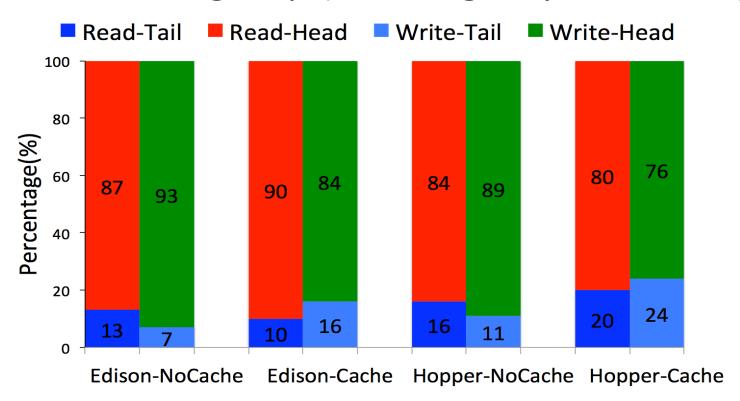
• Edison-NoCache, Write Response Time, 64MB



	Accuracy
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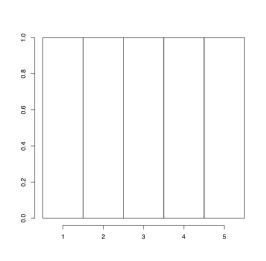
Percentage of Head group and Tail group

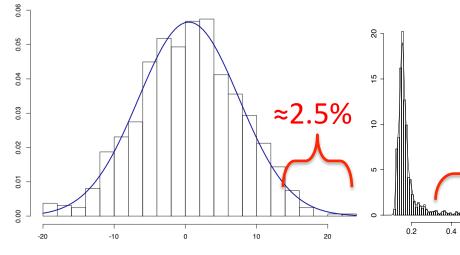
- 85% in Head group (i.e., small response time)
- 15% in Tail group (i.e., long response time)



What is Wrong with Using Normal or Uniform?

	Long Response Time (Rare Event)
Uniform Distribution	All equal
Normal Distribution	2.5%
Real Storage Systems (Edison and Hopper)	15%





Read (Stripe Size: 64MB)

≈15%

Response Time (sec.)

1.0

Summary

- Distribution function of response time of storage system is essential in estimating I/O performance
- We collected 2,537,900 response time sampling from 48 OSTes of 2 petascale storage systems across 141 days
- We found that single Normal or single Power law does not fit the response time
- We found that "Normal + Power law" fits response time better
- Future work
 - sample other storage systems
 - build accurate performance model
 - apply model to applications

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 Office of Advanced Scientific Computing Research, Office of Science, U.S. Department of Energy (Program manager: Lucy Nowell), support for the SDS project under contract number DE-AC02-05CH11231



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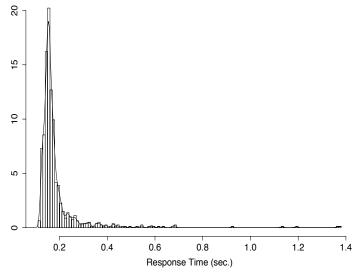
Thanks, Questions?

other questions, please email to: dbin@lbl.gov

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