

Fault-Tolerant Deep Learning Cache with Hash Ring for Load Balancing in HPC Systems

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Distributed Deep Learning in HPC

- The Data Parallel Approach
 - Replicates the deep learning model across nodes while distributing the training dataset among all nodes. *Parameter update*





Distributed Deep Learning in HPC

- Three key aspects of Distributed Deep Learning
 - I/O, Computation, Communication
- Most prior research efforts have concentrated on improving computation and communication.





Distributed Deep Learning in HPC

- However, the optimizations in computation and communication, along with the development of modern computational accelerators and network technologies, have shifted the bottleneck towards I/O.
- I/O accounts for 67-85% of total training time^[1].
 - Training ResNet50 on ImageNet: 85% of training runtime is IO overhead



[1] N. Dryden, R. Böhringer, T. Ben-Nun, and T. Hoefler, "Clairvoyant Prefetching for Distributed Machine Learning I/O," Proceedings of the International Conference for High Performance Computing, Networking, Storage and Analysis (SC '21), 2021.



Characteristics of Deep Learning Image Dataset

"Large Number of Small Files"

- ImageNet-1K: 1.28 million images (-150KB)
- ImageNet-21K: 11 million images (-163KB)
- **OpenImages:** 9 million images (-150KB)
- Google Landmarks Dataset v2: 5 million images (-200KB)
- Places365: 10 million images (-150KB)
- The HPC I/O subsystem is *not designed to efficiently handle the large-scale data I/O access required by deep learning frameworks.*



Optimizing I/O for Deep Learning Workloads

- **NoPFS**^[1]
 - Optimizes prefetching and caching by predicting data access patterns, reducing latency in training I/O.
- **DeepIO**^[2]
 - Minimizes backend storage reads by keeping data in memory, focusing on reducing read latency and boosting I/O efficiency for distributed training.
- HVAC^[3]
 - Caches data on node-local NVMe, specifically reducing repetitive I/O reads during training.

[2] Y. Zhu, "Entropy-Aware I/O Pipelining for Large-Scale Deep Learning on HPC Systems," Proceedings of the IEEE 26th International Symposium on Modeling, Analysis, and Simulation of Computer and Telecommunication Systems (MASCOTS), 2018.
 [3] A. Khan, A. K. Paul, C. Zimmer, S. Oral, S. Dash, S. Atchley, and F. Wang, "HVAC: Removing I/O Bottleneck for Large-Scale Deep Learning Applications," Proceedings of the IEEE International Conference on Cluster Computing (Cluster), 2022.



HVAC: High-Velocity AI Cache

- *HVAC*, a transparent read-only caching layer for large-scale supercomputers using node-local NVMe.
- Scalability
 - Designed to scale across thousands of compute nodes on leadership-class supercomputers like Summit and Frontier.
 - Avoids additional metadata bottlenecks and storage overhead.

Client-Server Library Architecture

- Intercepts <open-read-close> file I/O operations via LD_PRELOAD using a shared library approach.
- Data is cached to distributed node-local storage.
- Utilizes distributed hashing to determine the location of cached content across nodes. \rightarrow *No Repeated Access to PFS*



HVAC Overview

• HVAC Server

- Builds a caching layer on *node-local fast storage.*
- Handles system calls forwarded by HVAC clients.
- Reads files from node-local storage if available, or retrieves files from the PFS and caches them to node-local storage.

• HVAC Client

 Intercepts system calls directed to the PFS and redirects them to the HVAC server.







Increasing Node Failures in HPC and AI Workloads

- Node Failure Rates Rise with Complexity
 - Larger, more complex HPC systems have higher node failure rates.
 - Failure rates scale with system size, increasing linearly as more processors are added^[4].



• Intensive Workloads Increase Failure Probability

- High-demand tasks, such as large-scale deep learning jobs, also increase the risk of failure.
- Running multiple nodes for deep learning exacerbates the likelihood of failure event.





Increasing Node Failures in HPC and AI Workloads

Node Failure Rates Rise with Complexity



However, HVAC currently *lacks fault tolerance support*, which limits its resilience in handling node failures.

the likelihood of failure event.





Failures in HVAC

- Even a single node failure can halt the entire training process, despite fault-tolerance support in the DL framework.
 - E.g. Elastic Scaling *Horovod Elastic Run, MPI ULFM*
- This happens because I/O flows are controlled by HVAC.

\rightarrow The job must be restarted.



Failures in HVAC

• Even a single node failure can halt the entire training process,

Therefore, it is crucial to *ensure fault tolerance* in the *HVAC layer* to prevent training interruptions!



Naïve Approach

- Because HVAC functions as a caching layer, the original data resides in the PFS.
- I/O requests to failed nodes \rightarrow redirected to PFS!





I/O Redirection to PFS





Limitations of PFS I/O Redirection

"Frequent future PFS access"

- 22 mins per epoch x 5 epochs → *nearly 2 hours*
- Partial data reads from PFS still cause delays





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- Job Time Limitations: Risk of exceeding pre-defined job time.





Limitations of PFS I/O Redirection

"Re-caching Mechanism"

Read from PFS once, then access re-cached data for future requests

• Job Time Limitations: Risk of exceeding pre-defined job time.



Challenge: Handling Data Redistribution

- Original HVAC Implementation Static Hash Partitioning
 - Data paths are converted to key values and distributed across nodes using a modulo operation.
 - On node failure, recalculating hash values for N-1 nodes causes extensive data redistribution.





Challenge: Handling Data Redistribution

Additional Hash Functions

• Reduces data movement but doesn't address multiple unpredictable failures.

Range Partitioning

• Can handle multiple node failures, but *balancing data distribution* remains challenging.





Problem Definition

- How can we *track data locations* that change after re-caching?
- How can we *redistribute* lost data *evenly* across remaining active nodes?





Design & Implementation



Design of FT-HVAC

• **FT-HVAC:** An **I/O accelerated caching framework** with **fault tolerance** for large-scale distributed deep learning.



Design of FT-HVAC

- **FT-HVAC:** An **I/O accelerated caching framework** with **fault tolerance** for large-scale distributed deep learning.
- 1. Enable *fault tolerance* in HVAC.
- 2. Implement *data recaching* within the HVAC layer to ensure data availability and quick access during node failures.
- 3. Achieve *load-balanced* data recaching.



Elastic Recaching with Hash Ring

• Hash Ring Mechanism



<Before Failure>

Node Table	Hash Table	
Node 0 🛑 0.000 0.500 Node 1 🔵 0.375 0.875	File A	0.083427 •
Node 2 0.125 0.625 Virtual nodes	File E	0.293853



Elastic Recaching with Hash Ring

• Hash Ring Mechanism



<Before Failure>





Elastic Recaching with Hash Ring

• Hash Ring Mechanism



<Before Failure>



Elastic Recaching with Hash Ring

• Hash Ring Mechanism



<After Failure>















































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Evaluation



Node Failure Analysis on Frontier

- Analyzed job scheduler logs for six months following the production launch of the ORNL's Frontier cluster.
- Focused on three types of job failures: "Job Fail", "Node Fail", and "Timeout".

Туре	Count	Failure ratio	Overall ratio
Total Jobs	181,933	N/A	100%
Total Failures	45,556	100%	25.04%
Node Fail	1,174	2.58%	0.65%
Timeout	20,464	44.92%	11.25%
Job Fail	23,918	52.50%	13.15%

- **Job Fail:** Due to code errors, data/environment issues, or external malfunctions.
- Node Fail: Caused by hardware, network, software bugs, or overload.
- *Timeout:* Job exceeded set time limit, often due to complexity or resource/network constraints.



Node Failure Analysis on Frontier

• Average Runtime of Failed Jobs on Frontier



- Failed jobs typically run for an average of **over 1 hour**, sometimes reaching **2-3 hours.**
- Long-running job failures → significant loss of computing resources and time.
- Job failures have occurred consistently on a weekly basis.



Node Failure Analysis on Frontier

• Relationships between *types of job failures* and system variables.



(a) As the **number of nodes increases**, the rate of "**Node Failures**" **also rises**.

- E.g. With 7,750–9,300 nodes, "Node Failures" are **46.04%** of failures; including "Timeouts," they total 78.60%.
- (b) Execution time does not significantly impact the proportion of failure types.



Node Failure Analysis on Frontier

• Relationships between *types of job failures* and system variables.



(a) As the **number of nodes increases**, the rate of "**Node Failures**" **also rises**.

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Failures are *highly frequent* in large-scale HPC systems,

Software systems that are not designed to handle such failures are especially vulnerable, often resulting in significant, unavoidable losses

Experimental Setup

Attribute	Description
Supercomputer	Frontier
CPU	AMD Trento EPYC 7A53
GPU	8 x MI250X AMD with 64 GiB HBM
Memory Capacity	512 GiB DDR4
Node-local Storage	2 x 1.9 TB Samsung PM9A3 M.2 NVMe



- Application: Cosmoflow- MLPerf HPC v0.5 benchmark
- *Framework:* Horovod Elastic Run
- Dataset: 1.3TB cosmoUniverse dataset from NERSC ExaLearn group (524,288 training samples, 65,536 validation samples).
- File System: Orion (Lustre).
- *Training Setup:* 5 epochs, with 5 random failures injected after the first epoch.



Overhead Analysis

• Comparison between the original HVAC (NoFT) and two fault-tolerant approaches (FT w/PFS, FT w/NVMe) without any failure events.



- Overhead was minimal, with a maximum of 1-minute increase.
- Overhead resulted from additional data structures and conditional checks for the fault detection algorithm.



Overall Performance

• Performance evaluation with failure events.



- Comparison to No Failure Scenario
- FT w/PFS:
 - 64 nodes: 32.2% increase in training time.
 - 1024 nodes: 68.7% increase in training time.
- FT w/NVMe:
 - 64 nodes: 12.5% increase in training time.
 - 1024 nodes: 26.7% increase in training time.



Overall Performance

Performance evaluation with failure events.



- Comparison to No Failure Scenario
- FT w/PFS:
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 - 1024 nodes: 68.7% increase in training time.
- FT w/NVMe:
 - 64 nodes: 12.5% increase in training time.

FT w/ NVMe offers reduction in training time compared with FT w/ PFS, from **14.9%** to 24.9% as the number of nodes increases, as NVMe accesses the PFS only once following a failure, while FT with PFS requires repeated accesses.



Load Balance Analysis

• In a 1024-node configuration with **100 virtual nodes** per physical node, a single node failure redistributes data to an average of **80 nodes**.



Load Balance Analysis

• Analysis of varying virtual nodes per physical node to assess data distribution during failures.



- Simulation Setup: Conducted 500 simulations considering a 1024-node configuration.
- The Receiver Node metric represents the number of nodes receiving redistributed data.
- Increasing the number of virtual nodes improves data distribution but efficiency plateaus beyond 500.



Conclusion

- Node failures are *common* in leading-edge supercomputers.
- Such failures pose a *high risk to DL applications* on large-scale systems.
- FT-HVAC is a *fault-tolerant, I/O-accelerated caching framework* for distributed DL.
- FT-HVAC has demonstrated effective fault handling across 1024 nodes.
- The Elastic Recaching approach reduced training time by up to 24.9% compared to the I/O redirection method, while maintaining effective load balancing.



Questions?

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<Camera-ready paper>